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Edward C. Nevill

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EXAMINER

LI, AIMEE J

ART UNIT

PAPER NUMBER

2183

DATE MAILED: 01/12/2005

Please find below and/or attached an Office communication concerning this application or proceeding.

## Office Action Summary

Application No.

09/887,559

Applicant(s)

NEVILL ET AL.

Examiner

Aimee J Li

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-- The MAILING DATE of this communication appears on the cover sheet with the correspondence address --

### Period for Reply

A SHORTENED STATUTORY PERIOD FOR REPLY IS SET TO EXPIRE 3 MONTH(S) FROM THE MAILING DATE OF THIS COMMUNICATION.

- Extensions of time may be available under the provisions of 37 CFR 1.136(a). In no event, however, may a reply be timely filed after SIX (6) MONTHS from the mailing date of this communication.
- If the period for reply specified above is less than thirty (30) days, a reply within the statutory minimum of thirty (30) days will be considered timely.
- If NO period for reply is specified above, the maximum statutory period will apply and will expire SIX (6) MONTHS from the mailing date of this communication.
- Failure to reply within the set or extended period for reply will, by statute, cause the application to become ABANDONED (35 U.S.C. § 133). Any reply received by the Office later than three months after the mailing date of this communication, even if timely filed, may reduce any earned patent term adjustment. See 37 CFR 1.704(b).

### Status

- 1) ☒ Responsive to communication(s) filed on 30 September 2004 and 03 November 2004.
- 2a) ☒ This action is **FINAL**. 2b) ☐ This action is non-final.
- 3) ☐ Since this application is in condition for allowance except for formal matters, prosecution as to the merits is closed in accordance with the practice under *Ex parte Quayle*, 1935 C.D. 11, 453 O.G. 213.

### Disposition of Claims

- 4) ☒ Claim(s) 1-17 is/are pending in the application.
- 4a) Of the above claim(s) \_\_\_\_\_ is/are withdrawn from consideration.
- 5) ☐ Claim(s) \_\_\_\_\_ is/are allowed.
- 6) ☒ Claim(s) 1-17 is/are rejected.
- 7) ☐ Claim(s) \_\_\_\_\_ is/are objected to.
- 8) ☐ Claim(s) \_\_\_\_\_ are subject to restriction and/or election requirement.

### Application Papers

- 9) ☐ The specification is objected to by the Examiner.
- 10) ☐ The drawing(s) filed on \_\_\_\_\_ is/are: a) ☐ accepted or b) ☐ objected to by the Examiner.  
Applicant may not request that any objection to the drawing(s) be held in abeyance. See 37 CFR 1.85(a).  
Replacement drawing sheet(s) including the correction is required if the drawing(s) is objected to. See 37 CFR 1.121(d).
- 11) ☐ The oath or declaration is objected to by the Examiner. Note the attached Office Action or form PTO-152.

### Priority under 35 U.S.C. § 119

- 12) ☐ Acknowledgment is made of a claim for foreign priority under 35 U.S.C. § 119(a)-(d) or (f).
- a) ☐ All b) ☐ Some \* c) ☐ None of:
- ☐ Certified copies of the priority documents have been received.
  - ☐ Certified copies of the priority documents have been received in Application No. \_\_\_\_\_.
  - ☐ Copies of the certified copies of the priority documents have been received in this National Stage application from the International Bureau (PCT Rule 17.2(a)).
- \* See the attached detailed Office action for a list of the certified copies not received.

### Attachment(s)

- ☒ Notice of References Cited (PTO-892)
- ☐ Notice of Draftsperson's Patent Drawing Review (PTO-948)
- ☒ Information Disclosure Statement(s) (PTO-1449 or PTO/SB/08)  
Paper No(s)/Mail Date 03 November 2004.
- ☐ Interview Summary (PTO-413)  
Paper No(s)/Mail Date. \_\_\_\_\_.
- ☐ Notice of Informal Patent Application (PTO-152)
- ☐ Other: \_\_\_\_\_.

### DETAILED ACTION

1. Claims 1-17 have been examined. Claims 1-17 have been amended as per Applicant's request.

#### *Papers Submitted*

2. It is hereby acknowledged that the following papers have been received and placed on record in the file: Amendment as filed 05 October 2004 and IDS as filed 03 November 2004.

#### *Claim Rejections - 35 USC § 102*

3. The following is a quotation of the appropriate paragraphs of 35 U.S.C. 102 that form the basis for the rejections under this section made in this Office action:

A person shall be entitled to a patent unless –

(b) the invention was patented or described in a printed publication in this or a foreign country or in public use or on sale in this country, more than one year prior to the date of application for patent in the United States.

4. Claims 1-7, 11-12 and 16-17 are rejected under 35 U.S.C. 102(b) as being anticipated by Robinson et al., U.S. Patent No. 5,307,504.

5. Regarding claim 1, Robinson has taught an apparatus for processing data, said apparatus comprising:

- a. A processor core (20 of Fig.1) operable to execute operations as specified by instructions of a first instruction set (see Col.5 lines 51-55),
- b. An instruction translator (40 of Fig.2) operable to translate instructions of a second instruction set into translator output signals corresponding to instructions of said first instruction set (see Col.6 lines 20-43), at least one instruction of said second instruction set specifying an operation to be executed using one or more input variables (see Col.6 lines 52-57),

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- c. An interrupt handler responsive to an interrupt signal to interrupt execution of operations corresponding to instructions of said first instruction set after completion of execution of any currently executing operation (see Col.7 line 61 – Col.8 line 2),
  - d. Restart logic for restarting execution after said interrupt (see Col.7 lines 38-60).
  - e. Wherein said instruction translator is operable to generate a sequence of one or more sets of translator output signals corresponding to instructions of said first instruction set to represent said at least one instruction of said second instruction set (see Col.7 lines 6-37), each sequence being such that no change is made to said one or more input variables until a final operation within said sequence is executed (see Col.11 lines 58-65),
  - f. After occurrence of an interrupt during execution of a sequence of operations representing said at least one instruction of said second instruction set:
    - i. If said interrupt occurred prior to starting execution of a final operation in said sequence, then said restart logic restarts execution at a first operation in said sequence (see Col.7 lines 38-60),
    - ii. If said interrupt occurred after starting execution of a final operation in said sequence, then said restart logic restarts execution at a next instruction following said sequence (see Col.7 line 61 – Col.8 line 33).
6. Regarding claim 2, Robinson has taught an apparatus as claimed in claim 1, wherein said translator output signals include signals forming an instruction of said first instruction set (see Col.6 lines 20-27).

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7. Regarding claim 3, Robinson has taught an apparatus as claimed in claim 1, wherein said translator output signals include control signals that control operation of said processor core and said control signals match control signals produced on decoding instructions of said first instruction set (see Col.6 lines 1-6). Here, the instructions (control signals) are the same for both the directly compiled RISC instructions and the translated CISC-to-RISC Instructions (see Fig.1).

8. Regarding claim 4, Robinson has taught an apparatus as claimed in claim 1, wherein said translator output signals include control signals that control operation of said processor core and specify parameters not specified by control signals produced on decoding instructions of said first instruction set (see Col.7 lines 6-37). Here, the instruction ordering criteria is an output of the X-Y translator to control the processor on how to further group the translated instructions. The ordering criteria are not used for the decoding of instructions already in the Y (native) format (see Fig.1).

9. Regarding claim 5, Robinson has taught an apparatus as claimed in claim 1, wherein said restart logic is part of said instruction translator (see Col.7 line 3 – Col.8 line 33). Here, the method of interrupt handling and restarting of execution following the interrupt is performed within the translation system (40 of Fig.2), and thus the logic to perform the interrupt handling and restarting of execution is inherently within the translation system as well.

10. Regarding claim 6, Robinson has taught an apparatus as claimed in claim 1, wherein said restart logic stores a pointer to a restart location within instructions of said second instruction set that are being translated, said pointer being advanced upon execution of said final operation (see Col.7 lines 38-50).

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11. Regarding claim 7, Robinson has taught an apparatus as claimed in claim 6, wherein said pointer is a program counter value pointing to a memory address of a memory location storing an instruction of said second instruction set currently being translated (see Col.7 lines 38-50).

12. Regarding claim 11, Robinson has taught an apparatus as claimed in claim 1, wherein said input variables include system state variables not specified within said at least one of said second instruction set (see Col.7 lines 15-21). Here, translated Y instructions specify temporary variables and storage locations prior to updating the state variables and storage locations specified by the non-translated X instructions.

13. Regarding claim 12, Robinson has taught an apparatus as claimed in claim 1, wherein said processor has a register bank (97 of Fig.4) containing a plurality of registers and instructions of said first instruction set execute operations upon register operands held in said registers (see Col.7 lines 25-37).

14. Regarding claim 16, Robinson has taught a method of processing data, said method comprising the steps of:

- a. Executing operations as specified by instructions of a first instruction set (see Col.5 lines 51-55).
- b. Translating instructions of a second instruction set into translator output signals corresponding to instructions of said first instruction set (see Col.6 lines 20-43), at least one instruction of said second instructions et specifying an operation to be executed using one or more input variables (see Col.6 lines 52-57),

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- c. In response to an interrupt signal, interrupting execution of operations corresponding to instructions of said first instruction set after completion of execution of any currently executing operation (see Col.7 line 61 – Col.8 line 2),
  - d. Restarting execution after said interrupt (see Col.7 lines 38-60),
  - e. Wherein said step of translating generates a sequence of one or more sets of translator output signals corresponding to instructions of said first instruction set to represent said at least one instruction of said second instruction set (see Col.7 lines 6-37), each sequence being such that no changes is made to said one or more input variables until a final operation within said sequence is executed (see Col.11 lines 58-65),
  - f. After occurrence of an interrupt during execution of a sequence of operations representing said at least one instruction of said second instruction set:
    - i. If said interrupt occurred prior to starting execution of a final operation in said sequence, then restarting execution at a first operation in said sequence (see Col.7 lines 38-60),
    - ii. If said interrupt occurred after starting execution of a final operation in said sequence, then restarting execution at a next instruction following said sequence (see Col.7 line 61 – Col.8 line 33).
15. Regarding claim 17, Robinson has taught a computer program product including a computer program for controlling a computer to perform the method of claim 16 (see Col.5 lines 43-50).

***Claim Rejections - 35 USC § 103***

16. The following is a quotation of 35 U.S.C. 103(a) which forms the basis for all obviousness rejections set forth in this Office action:

(a) A patent may not be obtained though the invention is not identically disclosed or described as set forth in section 102 of this title, if the differences between the subject matter sought to be patented and the prior art are such that the subject matter as a whole would have been obvious at the time the invention was made to a person having ordinary skill in the art to which said subject matter pertains. Patentability shall not be negated by the manner in which the invention was made.

17. Claims 8-10 and 13-15 are rejected under 35 U.S.C. 103(a) as being unpatentable over Robinson et al., U.S. Patent No. 5,307,504 as applied to claims 1 and 12 above, and further in view of Dickol et al., U.S. Patent No. 5,898,885.

18. Regarding claim 8, Robinson has taught an apparatus as claimed in claim 1, but has not explicitly taught wherein instructions of said second instruction set specify operations to be executed upon stack operands held in a stack and said input variables include input stack operands.

19. However, Dickol has taught non-native Java instructions that specify operations to be executed on operands held in a stack (see Dickol, Col.4 lines 42-61), and those instruction's corresponding translation into RISC-type native instructions so that redundant execution steps are eliminated and native code execution performance is improved (see Dickol, Col.2 lines 17-39). Because Robinson has taught an apparatus translating instructions from a non-native to a native instruction set (see Robinson, Col.5 lines 43-67), but hasn't explicitly specified if the instructions of the non-native instruction set specify operations to be executed on stack operands, one of ordinary skill in the art would have found it obvious to modify the translation apparatus of Robinson to instead translate instructions specifying operations to be performed on stack



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operands in a non-native second instruction set into a native first instruction set in a manner so that redundant execution steps are eliminated, thereby improving processor performance.

20. Regarding claim 9, Robinson in view of Dickol has taught an apparatus as claimed in claim 8, wherein any updating of the state (memory or registers) is not performed until after execution of said final operation has commenced in order to preserve instruction granularity (see Robinson, Col.7 line 38 - Col.8 line 34), but has not explicitly taught wherein the updating of the state includes removing stack operands from the stack.

21. However, Dickol has taught non-native Java instructions that specify operations to be executed on operands held in a stack (see Dickol, Col.4 lines 42-61), and the corresponding translation of those instructions into RISC-type native instructions so that redundant execution steps are eliminated and native code execution performance is improved (see Dickol, Col.2 lines 17-39). Because Robinson has taught an apparatus translating instructions from a non-native to a native instruction set (see Robinson, Col.5 lines 43-67) where updating of the state does not occur until after execution of a final operation has commenced (see Robinson, Col.7 line 38 - Col.8 line 34), and further because the a push or pop of Dickol is a state update (push/pop operates on the stack and registers), one of ordinary skill in the art would have found it obvious to modify the processor of Robinson to only update the state (i.e. pop registers from the stack) after execution of a final operation has commenced so that non-native to native instruction translation can be executed more efficiently (see Dickol, Col.7 lines 5-16) while maintaining instruction granularity (see Robinson, Col.7 lines 38-60).

22. Regarding claim 10, Robinson has taught an apparatus as claimed in claim 8, wherein any updating of the state (memory or registers) is not performed until after execution of said final

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operation has commenced in order to preserve instruction granularity (see Robinson, Col.7 line 38 - Col.8 line 34), but has not explicitly taught wherein the updating of the state includes adding stack operands to the stack.

23. However, Dickol has taught non-native Java instructions that specify operations to be executed on operands held in a stack (see Dickol, Col.4 lines 42-61), and the corresponding translation of those instructions into RISC-type native instructions so that redundant execution steps are eliminated and native code execution performance is improved (see Dickol, Col.2 lines 17-39). Because Robinson has taught an apparatus translating instructions from a non-native to a native instruction set (see Robinson, Col.5 lines 43-67) where updating of the state does not occur until after execution of a final operation has commenced (see Robinson, Col.7 line 38 - Col.8 line 34), and further because the a push or pop of Dickol is a state update (push/pop operates on the stack and registers), one of ordinary skill in the art would have found it obvious to modify the processor of Robinson to only update the state (i.e. push registers onto the stack) after execution of a final operation has commenced so that non-native to native instruction translation can be executed more efficiently (see Dickol, Col.7 lines 5-16) while maintaining instruction granularity (see Robinson, Col.7 lines 38-60).

24. Regarding claim 13, Robinson has taught an apparatus as claimed in claim 12, but has not explicitly taught wherein a set or registers within said register bank hold stack operands from a top portion of said stack.

25. However, Dickol has taught wherein a set of registers in a register file hold operands corresponding to a top portion of a stack (see Dickol, Col.2 lines 59-67). Here, when non-native Java paired push and pop instructions are to be translated into native RISC-type instructions, the

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translation stores those stack operands, which are on the “top” of the stack (due to the LIFO nature of stacks), in the register file, so that unnecessary data transfers are eliminated, and processor performance is improved (see Dickol, Col.4 lines 42-61). Because Robinson has taught an apparatus translating instructions from a non-native to a native instruction set (see Robinson, Col.5 lines 43-67) with the instructions able to operate on operands in the register file (see Robinson, Col.7 lines 25-37), but hasn’t explicitly specified that the operands in the register file are stack operands, one of ordinary skill in the art would have found it obvious to modify the translation apparatus of Robinson to instead translate instructions specifying operations to be performed on stack operands in a non-native second instruction set into a native first instruction set in a manner such that redundant stack operations have their operands stored in the register file so that redundant execution steps can be eliminated, thereby improving processor performance.

26. Regarding claim 14, Robinson has taught an apparatus as claimed in claim 13, wherein an apparatus translating instructions from a non-native to a native instruction set (see Robinson, Col.5 lines 43-67), where updating of the state does not occur until after execution of a final operation has commenced (see Robinson, Col.7 line 38 - Col.8 line 34), but has not explicitly taught wherein said instruction translator has a plurality of mapping states in which different registers within said set of registers hold respective stack operands from different positions within said stack, said instruction translator being operable to move between mapping states when said final operation is executed so as to update said input variables.

27. However, Dickol has taught mapping of registers to different stack operands with the ability to move between the mappings based on whether there is a push or a pop operation (see Dickol, Col.6 lines 27-40), thus allowing non-native to native instruction translation to be

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executed more efficiently (see Dickol, Col.7 lines 5-16). Because a push or pop instruction of Dickol is a non-native instruction (see Dickol, Col.6 lines 27-40), one of ordinary skill in the art would have found it obvious to modify the processor of Robinson to map registers of different stack operands with the ability to move between mappings once the final native operation of a non-native push or pop instruction has been executed so that non-native to native instruction translation can be executed more efficiently (see Dickol, Col.7 lines 5-16) while maintaining instruction granularity (see Robinson, Col.7 lines 38-60).

28. Regarding claim 15, Robinson has taught an apparatus as claimed in claim 1, but has not explicitly taught wherein said instructions of said second instruction set are Java Virtual Machine instructions.

29. However, Dickol has taught non-native Java Virtual Machine instructions being translated into RISC-type native instructions (see Dickol, Col.3 lines 49-57) so that redundant execution steps are eliminated and native code execution performance is improved (see Dickol, Col.2 lines 17-39). Because Robinson has taught an apparatus translating instructions from a non-native to a native instruction set (see Robinson, Col.5 lines 43-67), but hasn't explicitly specified the type of non-native instructions, one of ordinary skill in the art would have found it obvious to modify the translation apparatus of Robinson to instead translate Java Virtual Machine instructions into a native first instruction set in a manner so that redundant execution steps are eliminated, thereby improving processor performance.

### ***Response to Arguments***

30. Examiner withdraws objections to the Abstract and specification in favor of the amended specification and Abstract.

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31. Examiner withdraws the objection to claim 17 regarding failure to further limit in favor of the arguments presented by Applicant.

32. Examiner withdraws objections to claims 1-15 regarding minor informalities in favor of the amended claims.

33. Examiner withdraws the 35 U.S.C. § 112, second paragraph rejections of claims 3 and 11 in favor of the amended claims.

34. Applicant's arguments filed 05 October 2004 have been fully considered but they are not persuasive. Applicants argue in essence on pages 13-15

...Robinson fails to disclose that “no changes are made to the input variables until execution of a final operation in the sequence.”...

...

...Robinson teaches that changes are made to the input variables prior to execution of a final operation in the sequence...

35. This has not been found persuasive. Robinson teaches in column 11, lines 63-65 a “1- or 2-byte non-interlocked write on a single-processor machine,” whose translated sequence is a “read-modify-write sequence.” There is only one write in this sequence, which occurs in “the group 3 instruction.” In Robinson’s column 7, lines 6-37, there are four groups of instructions with the fourth group being another write in the sequence. In the situation recited in Robinson’s column 11, lines 63-65, the fourth group is not required, essentially becoming a no-operation. This means that the final operation in the sequence is the group 3 instruction, which writes back the state.

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36. In addition, the claim states “no change is made to said one or more input variables until a final operation within said sequence is executed.” Robinson has taught that the X state, which includes the X memory state and X register state, is updated in group 3 and group 4 instructions, but this does not mean that the input variables are changed. This just means that the memory and registers are updated with the results, i.e. something is written back to the X memory and X registers. The data that was operated on was held in the temporary storage, and “operated on” does not mean changed, only that some function or work was performed with the input. The results are initially put in temporary storage before being written back to permanent storage, but this does NOT mean that the input variables were changed in any fashion. The only way for the input to be changed is if the input permanent storage locations were specifically indicated as the destination in permanent storage for the results. However, Robinson does not state anything to this effect just that the results are written back to permanent storage.

### *Conclusion*

37. The prior art made of record and not relied upon is considered pertinent to applicant's disclosure as follows. Applicant is reminded that in amending in response to a rejection of claims, the patentable novelty must be clearly shown in view of the state of the art disclosed by the references cited and the objections made. Applicant must also show how the amendments avoid such references and objections. See 37 CFR § 1.111(c).

- a. Cohen et al., U.S. Patent Number 6,718,539, has taught an instruction translator that translates an instruction from one instruction set to a sequence of instructions from another instruction set.

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38. **THIS ACTION IS MADE FINAL.** Applicant is reminded of the extension of time policy as set forth in 37 CFR 1.136(a).

39. A shortened statutory period for reply to this final action is set to expire **THREE MONTHS** from the mailing date of this action. In the event a first reply is filed within **TWO MONTHS** of the mailing date of this final action and the advisory action is not mailed until after the end of the **THREE-MONTH** shortened statutory period, then the shortened statutory period will expire on the date the advisory action is mailed, and any extension fee pursuant to 37 CFR 1.136(a) will be calculated from the mailing date of the advisory action. In no event, however, will the statutory period for reply expire later than **SIX MONTHS** from the mailing date of this final action.

40. Any inquiry concerning this communication or earlier communications from the examiner should be directed to Aimee J Li whose telephone number is (571) 272-4169. The examiner can normally be reached on M-T 7:30am-5:00pm.

41. If attempts to reach the examiner by telephone are unsuccessful, the examiner's supervisor, Eddie Chan can be reached on (571) 272-4162. The fax phone number for the organization where this application or proceeding is assigned is 703-872-9306.

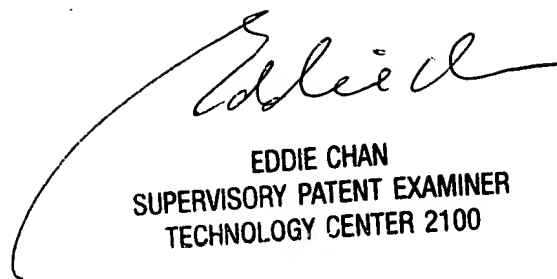
42. Information regarding the status of an application may be obtained from the Patent Application Information Retrieval (PAIR) system. Status information for published applications may be obtained from either Private PAIR or Public PAIR. Status information for unpublished applications is available through Private PAIR only. For more information about the PAIR system, see <http://pair-direct.uspto.gov>. Should you have questions on access to the Private PAIR system, contact the Electronic Business Center (EBC) at 866-217-9197 (toll-free).

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AJL

Aimee J. Li

07 January 2005



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